Hardening

Applied Information Security
Lecture 4

Recap: Foreknowledge

Attackers:

- Mindset, Phases

What can be accomplished by an attack:

- Code Injection
  Dynamic Evaluation, Insecure Deserialization, XSS, Command Inject, SQL Injection, Buffer Overflow, ...
- Side Channels
  Hardware, Network, Physical World (airgap), ...
- Social Engineering
  (spear)phishing, weapons of influence, dumpster diving, ...

Today's Topics

- detect vulnerabilities
  - known
    - vulnerability scan & exploit
  - unknown
    - automatic, manual
- detect attacks
  - ongoing
    - intrusion detection
  - past
    - auditing, malware removal
- administration
  - hardening, firewall, isolation

Detect Vulnerabilities: Known

Known vulnerabilities;
How to find them?
How to exploit them?

recently changed names; some still refer to this as **OpenVAS**

gvm

Greenbone Vulnerability Manager

MITRE CVE, + UI & auto-scan

- scan a specified host for vulnerabilities

Date	Status	Task	Severity	Scan Results				
				High	Medium	Low	Log	False Pos.
Thu Jan 9 03:05:08 2020	Done	Immediate scan of IP 192.168.11.137	N/A	0	0	0	0	

Report: Results (312 of 734)

ID: 97cc63d0-65d7-45ee-8ca8
Modified:
Created:
Owner: admin

1 - 10

Vulnerability	Severity	QoD	Host	Location
rexec Passwordless / Unencrypted Cleartext Login	10.0 (High)	75%	192.168.11.137	512/tcp
Samba End Of Life Detection	10.0 (High)	75%	192.168.11.137	445/tcp
Samba 'TALLOC_FREE()' Function Remote Code Execution Vulnerability	10.0 (High)	75%	192.168.11.137	445/tcp
PHP Multiple Vulnerabilities - Aug08	10.0 (High)	75%	192.168.11.137	80/tcp
PHP Version < 5.2.7 Multiple Vulnerabilities	10.0 (High)	75%	192.168.11.137	80/tcp
PHP End Of Life Detection (Linux)	10.0 (High)	75%	192.168.11.137	80/tcp
MySQL End Of Life Detection (Linux)	10.0 (High)	75%	192.168.11.137	3306/tcp
PostgreSQL End Of Life Detection (Linux)	10.0 (High)	75%	192.168.11.137	5432/tcp

msfconsole

Metasploit Framework Console

collection of exploits

- standard format (use, options, run)
- Parameterized (IP address, ...)

for security auditing

```
msf exploit(multi/handler) > run

[*] Started reverse TCP handler on 192.168.86.223:4444
[*] Sending stage (179779 bytes) to 192.168.86.61
[*] Meterpreter session 1 opened (192.168.86.223:4444 -> 192.168.86.61:49197) at
 2018-05-29 11:48:32 -0400

meterpreter > shell
Process 3028 created.
Channel 1 created.
Microsoft Windows [Version 6.1.7601]
Copyright (c) 2009 Microsoft Corporation.  All rights reserved.

C:\Users\victim\Downloads>
```

Detect Vulnerabilities: Unknown

No known vulnerabilities; How do we craft new ones?

- automatic
- manual

Automatic

afl, sqlmap

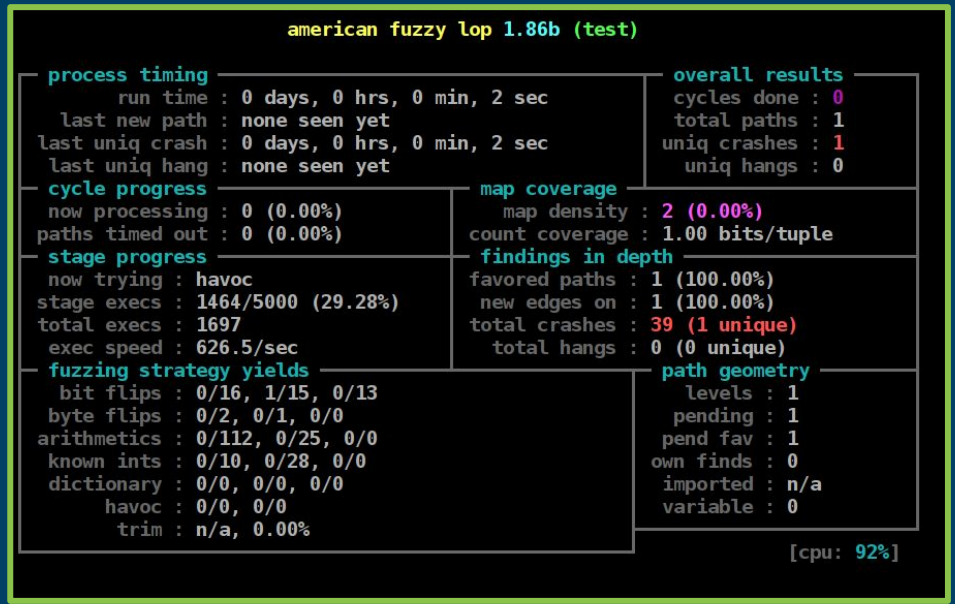
afl

American Fuzzy Lop

fuzzer

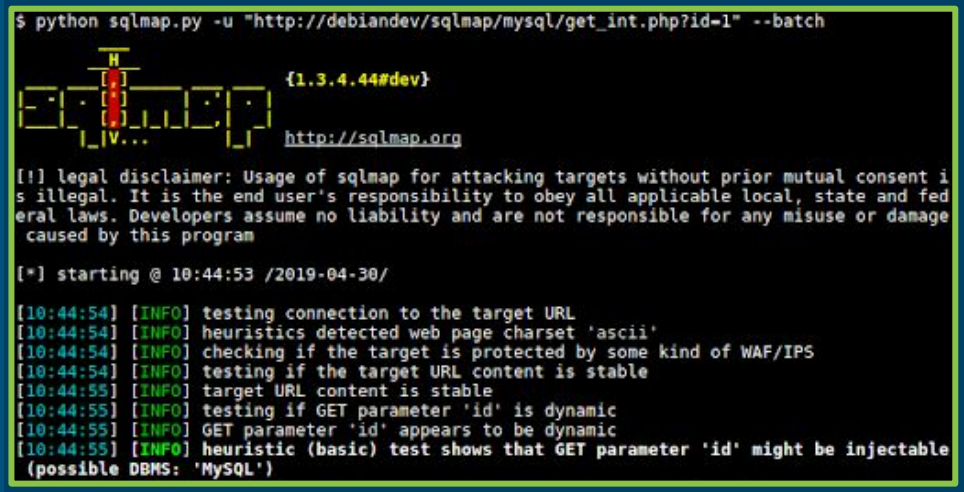
- randomly tweak a given input
- provoke bad behavior

used e.g. to find buffer overflows.



sqlmap

test for SQL injections!



Manual

angr, Ghidra (reverse engineer), Burpsuite, ...

angr

binary analysis: concolic execution

- symbolic execution
  (abstract execution traces)
- + testing
  (to enter each abstract trace)

used e.g. to find buffer overflows.

```
(angr) last@ubuntu: ~/angr_ctf/dist $ python scaffold00.py
WARNING | 2019-03-20 18:00:01,593 | angr.state_plugins.symbolic_memory | The program is accessing memory or regist
WARNING | 2019-03-20 18:00:01,593 | angr.state_plugins.symbolic_memory | angr will cope with this by generating ar
WARNING | 2019-03-20 18:00:01,593 | angr.state_plugins.symbolic_memory | 1) setting a value to the initial state
WARNING | 2019-03-20 18:00:01,594 | angr.state_plugins.symbolic_memory | 2) adding the state option ZERO_FILL_UNC
WARNING | 2019-03-20 18:00:01,594 | angr.state_plugins.symbolic_memory | 3) adding the state option SYMBOL_FILL_U
WARNING | 2019-03-20 18:00:01,594 | angr.state_plugins.symbolic_memory | Filling register edi with 4 unconstrained
WARNING | 2019-03-20 18:00:01,598 | angr.state_plugins.symbolic_memory | Filling register ebx with 4 unconstrained
WARNING | 2019-03-20 18:00:03,181 | angr.state_plugins.symbolic_memory | Filling memory at 0x7fff0000 with 83 unco
WARNING | 2019-03-20 18:00:03,182 | angr.state_plugins.symbolic_memory | Filling memory at 0x7ffeff60 with 4 uncon
[+] Success! Solution is: JXWVXRKX
(angr) last@ubuntu: ~/angr_ctf/dist $ ./00_angr_find
Enter the password: JXWVXRKX
Good Job.
(angr) last@ubuntu: ~/angr_ctf/dist $
```

Detect Attacks: Ongoing

Attack is taking place!
How will we know?
How will we handle it?

- Intrusion detection

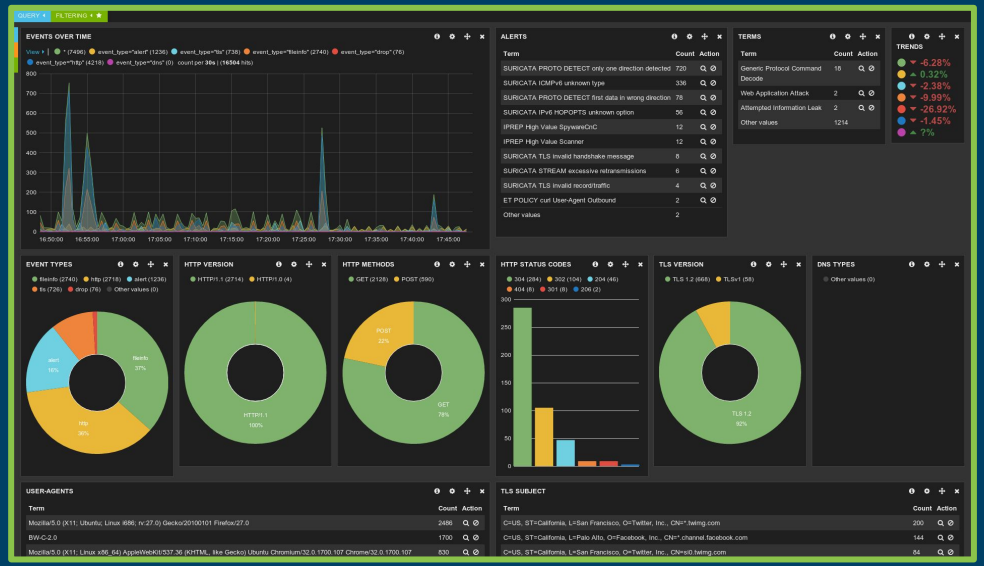
Intrusion Detection

Intrusion Detection System

- automated review and response
- responds in (nearly) real time
- components:
  - sensors
  - analysis engine
  - countermeasure deployment
  - audit log

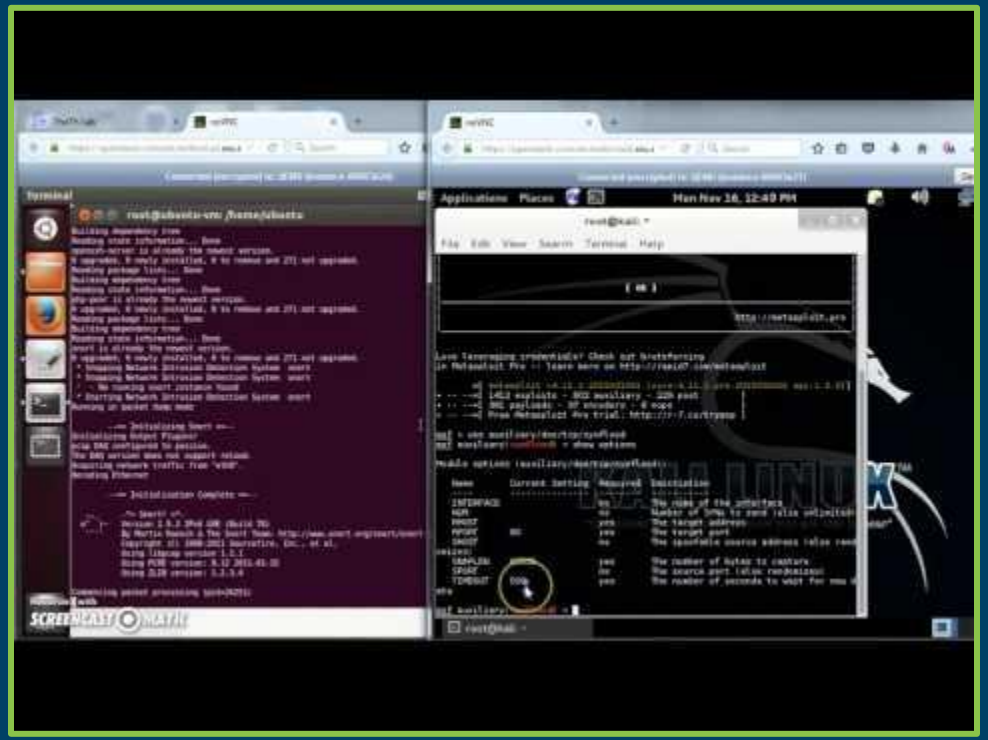
Example: Network Monitoring

- suspicious behavior:
  - opening connections to many hosts
- automated response:
  - router reconfigures to isolate suspicious host on its own subnet with access only to (e.g.) virus scanner download.
  - notifies administrators
- issue:
  - errors...



Screenshot from “Suricata” (Network Intrusion Detection and Prevention tool)

Example: Intrusion Detection, Snort



Network-Based Intrusion Detection System

typically separate machine

- **stealth mode:**
  - one NIC faces the network being monitored,
    no packets ever sent out on it,
    no packets can be routed specifically to it
  - another NIC faces a separate network
    through which alarms are sent
- **honeypot:**
  - dedicated machines(s) or networks
  - purpose is to look attractive to attacker
  - but actually just a trap: monitored to detect and surveil attacker

Errors

- **false positive**
  raise an alarm for a non-attack
  - makes administrators less confident in warnings
  - perhaps leading to actual attacks being dismissed
- **false negative**
  not raise an alarm for an attack
  - the attackers get in undetected!
- tradeoff between the two needs to be tunable;
  difficult to achieve the right classification statistics

(problem if both possible at the same time. cf. type soundness)

Identification Methodologies

[Denning 1987]

- signature based:
  - recognize known attacks
- specification based:
  - recognize bad behavior
- anomaly based:
  - recognize abnormal behavior

Signature Based Detection

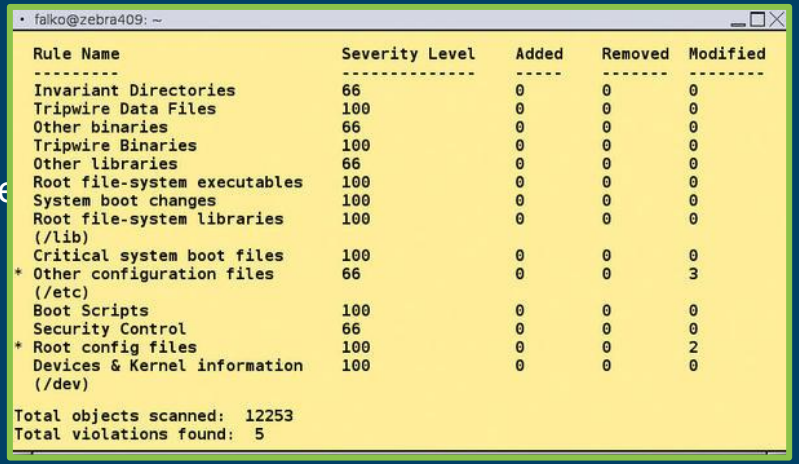
a.k.a. {misuse, rule-based} detection

- characterize known attacks w/ signatures
- behavior matches signature ⇒ declare an intrusion
- issues:
  - works only for known attacks
  - signature needs to be robust w.r.t. small changes in attack

Example: Tripwire

open source tool and commercial product

- **policy**:
  - certain files shouldn't change
- **state snapshot**:
  - analyzes filesystem, stores database of file hashe
- **automated response**:
  - runs (e.g. daily) and reports change of hash
- **issues**:
  - where to store database, how to protect its integrity, how to protect tripwire itself?



Example: Network Flight Recorder

Network Flight Recorder

- three components:
  - packet sucker captures network traffic
  - decision engine uses custom-written filters in DSL to extract information from packets
  - backend writes information to disk; packets are discarded
- queries performed over stored information
  while rest of system continues to process packets
- similar ideas used in **Zeek** (aka. **Bro**)
  [Paxson 1999], available still as open source IDS

Specification Based Detection

$$
\begin{aligned}
&\text{Rule} \equiv \text{deniedRule} \sqcup \text{permittedRule} \\
&\text{deniedRule} \sqsubseteq \neg\text{permittedRule} \\
&\text{Rule} \equiv \exists\ \text{hasSrc} \sqcap \forall\text{hasSrc.NetworkRole} \sqcap \\
&\qquad \exists\ \text{hasDst} \sqcap \forall\text{hasDst.NetworkRole} \sqcap \\
&\qquad \exists\ \text{hasService} \sqcap \forall\text{hasService.Service} \sqcap \\
&\qquad \exists\ \text{hasOrder} \sqcap \exists\ \text{hasDir} \sqcap \forall\text{hasDir.Direction} \\
&\text{Direction} \equiv \text{inDir} \sqcup \text{outDir} \qquad \text{inDir} \sqsubseteq \neg\text{outDir}
\end{aligned}
$$

- characterize good behavior of program w/ a specification
- if behavior ever departs from specification, declare an intrusion
- issues:
  - effort to create specifications
  - any program is a potential vulnerability if executed by a privileged user

Example: Distributed Monitor

[Ko et al. 1997]

- monitors Unix audit logs
- analyst writes grammar in DSL to describe good behavior
- parser checks conformance of logs with grammar
- distributed because it combines information from multiple hosts

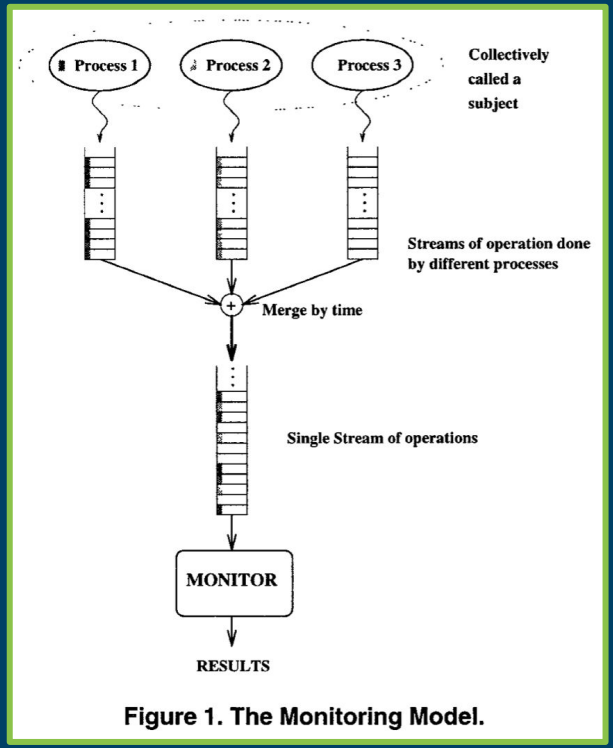
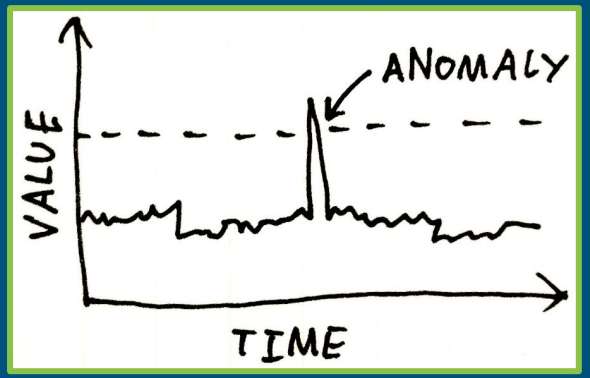


Figure 1. The Monitoring Model.

Anomaly Based Detection



- characterize normal behavior of system
- if behavior ever departs far enough from normal, declare an intrusion
- issues:
  - feature identification
  - obtaining data on what is normal →

Example: Haystack (US Air Force)

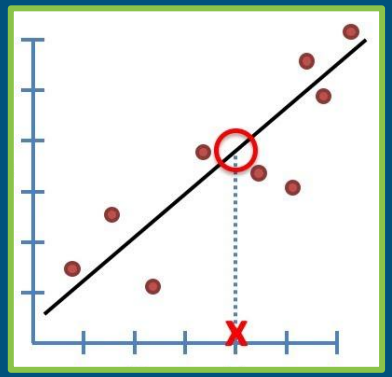
[Smaha 1988] (influential; one of the earliest IDS papers)

- monitors value of some statistic of interest over a sliding time window: a_i , a_{i+1}, ..., a_j ← (buffer j-i msg)
- determine lower and upper bounds t_L and t_U such that 90% of values lie between t_L and t_U
- next value is outside t_L and t_U ⇒ anomaly; raise alarm

adaptive

- as time passes, window moves, so detector adjusts itself.

Statistical Models



ML great for *classification*...

- **threshold models**
  - min and max
- **moment models**
  - mean & standard deviation
- **markov models**
  - probability of next event based on current state

BUT, ML **not** great for *outlier detection*. Adversarial ML poorly understood.

Intrusion Response?

intrusion handling: [Northcutt 1998]

1. Preparation
2. Identification
3. Containment
4. Eradication
5. Recovery
6. Follow up

automated response: monitor, protect, alert

counterattack?

- **legal route:** file criminal complaint
- tech route: damage attacker
  - might harm innocents
  - might expose you legally

Detect Attacks: Past

Damage is done.
How will we know?
How will we handle it?

- Auditing
- Anti-malware

Auditing

What to Log?

example: US State Dept. pilot program (1980s)

- requirements
  - log every transaction related to protected electronic documents
  - system administrator reviews log daily to search for malicious behavior
- experiment
  - test system for 5 users, 10 minutes
- result
  - audit log = stack of paper over 1ft high
  - real system would have been 1000s of users working 24/7
- lessons learned
  - logging and review of everything by a human is impractical
  - need: reduce information logged
  - need: automated review

States vs. Events

States

data; what system *is*

- backup
- more?

what state to log?

pros: survive power failures, crashes, attacks.
cons: what state? memory, disk, network, …
what about distributed systems? (hard)

Events

our focus

actions, how system *came to be*

- login
- access to protected resource,
- elevation and attenuation of privileges,
- ...

which events to log?

- event relevant for security
- what check was made, outcome, information that lead to that decision.

In-Class: Course Management System

what kind of events to log for a course management system (*mutations*)?
what details would you put into the log entry?

Log Type	Action	Acting NetID	Acting IP Address	Affected NetIDs	Simulated NetID	Assignment	Date
Course	Created New Assignment	mrc26	128.84.217.18			A1, Homework 4	January 28, 2016 04:06PM

- Created new assignment 'A1'
- Added required submission 'a1' with accepted types: pdf
- Added problem 'a1' worth 4.0 points
- Created new groups for each student

Auditing - What to Log

In-Class: Course Management System

Course

- add students,
- change group timeslot,
- computed assignment stats,
- computed total score,
- created / edited removed / restored assignment,
- created / removed / restored announcement,
- created / removed timeslot,
- dropped students,
- edited course properties,
- edited staff preferences,
- edited student preferences,
- sent course email
- uploaded class list

Content

- added / edited content data,
- create / edited / reorder / remove content,
- add students,
- change group timeslot,

Group

- sent / canceled group invite
- joined / left group
- created / disbanded group
- granted / removed extension
- requested regrade,

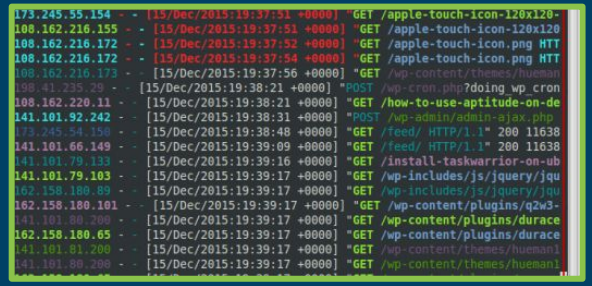
Grade

- assigned grader
- edited grades
- edited comments
- uploaded grade files

In-Class: Course Management System

details logged:

- event type
- acting NetID
- acting IP address
- affected NetIDs
- simulated NetID
- assignment, if any
- event details (no sanitization of grades)



How to Log

Say what you mean

log entry should **say what it means**.

- interpretation of log entry should depend only on content of log entry
- ⇒ reviewer can recover meaning w/o needing to assume / supply context
- good practice: write down straightforward English sentence describing the meaning of each log entry

Standard Log File Format

keeping log files in standard format enables...

- reuse of tools for log analysis
- correlation across logs from multiple applications

standard formats:

- Common Log Format (NCSA; used by web servers)
- **syslog** (used by Unix)
  - originated with sendmail
  - became de facto standard
  - then standardized by IETF: RFC 5424
  - examples: take a look in your local /var/log directory

Log Size Too Large?

what happens if log size grows too large?

- stop logging
- overwrite previous entries
- halt system

all used in practise, depending on scenario.

(none of these options are great. but you have to do something)

Manual

enable admins to explore logs and look for {states,events}.

issues:

- designers might not have recorded the right {states,events}
- visualization, query, expressivity (HCI/DB issues)
- correlation amongst multiple logs

Auditing - Manual

Visualization

Interface

- text
  - example: syslog (previous slide)
- hypertext
- DBMS
  - example: queries in the course management system
- graph
  - nodes might be entities (processes, files), edges might be associations (forking, times)

Available Pages

High Level Pages

- date page

Auid (User) Pages

- **jhoaglan**

Process Info Pages

Parentless Processes (creating fork not in log)

- **115**
- **499**

Parented Processes

- **500**: /usr/bin/csh
- **501**: /usr/ucb/quota
- **502**: /usr/bin/tty
- **505**: /usr/bin/ln
- **506**: /usr/bin/sh
- **507**: /usr/bin/id
- **508**: /usr/bin/chmod

File Info Pages

- **/etc/security/password.adjunct**
- **/usr/export/home/jhoaglan/.history**

Generated on Wed Jul 20 12:05:28 PDT 1994 by hoagland using ab

Figure 3d. Output of the Hypertext Generator on the log of a suid attack: the index page.

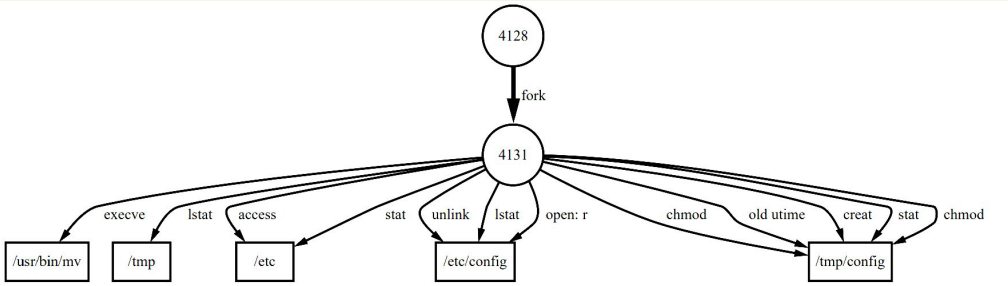


Figure 4b. FAB presentation of "Investigating the Disappearance of a File" situation with focus on process 4131. The process that removed the file, 4131, is now investigated further with other system calls and its parent process shown. It is seen that process is running the "mv" program, which is show to read and remove "/etc/config",

Auditing - Manual

Visualization

Interface

- text
  - example: syslog (previous slide)
- hypertext
- DBMS
  - example: queries in the course management system
- graph
  - nodes might be entities (processes, files),
    edges might be associations (forking, times)

Technique

- temporal
  - animate what happened and when (e.g. time-ordered sequence of graphs)
- slice
  - minimal set of log events that affected an object

Automatic

example: LogRhythm

detect

- suspicious behavior
- violations of explicit policy

built how

- custom-built systems
- classic AI techniques like training neural nets, expert systems, etc.
- machine learning

response: monitor, report, take action

e.g. close account / connection

Malware removal

antivirus

Administration

What can Sys-Admins do, to secure systems, w/o writing them themselves?

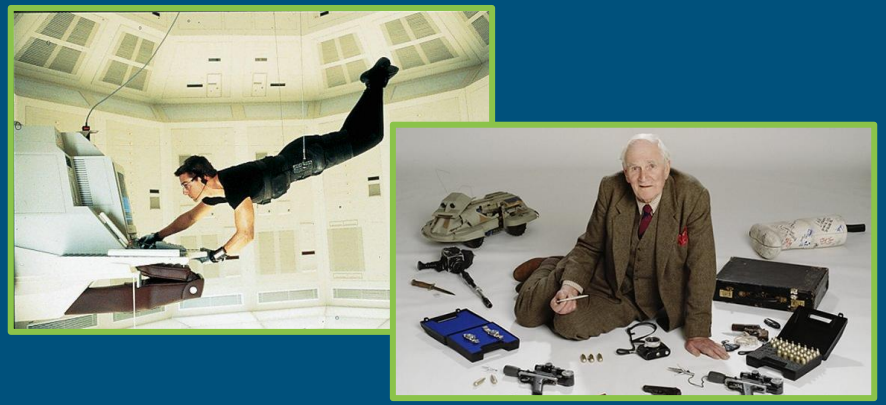
- Hardening
- Firewalls
- Isolation

Isolation Categories

Categories (Subcategories)	Trusted Computing Base	Performance Overhead of common cases*	Code Requirement
Physical Host	Hardware	Negligible	No code modification required
Hardware component	Hardware, isolation technique framework	High	Application source code may be required
Supervisor			
Hypervisor	Hardware, isolation technique framework, or BIOS	Practical	Application source code or binary code may or may not require to be modified
Library OS	Hypervisor, OS, Library OS framework	High	Application code may require to be ported or recompiled
Container	OS, sandbox framework, Container Engine, Piece of Application Logic	Practical	No code modification required
Intra-application			
Code Rewriting	OS and Binary Writer	Practical	Modification to binary required.
Compiler	Operating system, compiler and runtime	Low	Source code modification may or may not be required
System Loading	OS kernel, isolation technique framework	High	Application code modification may or may not be required.

isolation done where:
outside the computer

Physical

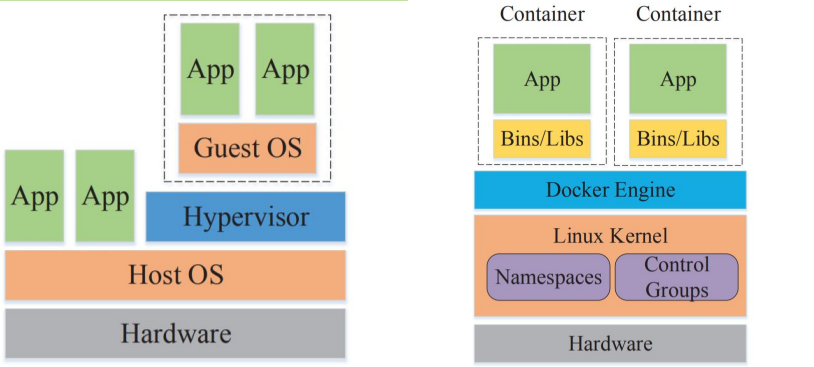


air-gap: system physically isolated from network.

how to breach (w/ enough **resources**)? Hollywood-level creative.

isolation done where:
outside the program

Supervisor



(b) Type-2 Hypervisor

Fig. 5: Architectural Overview of the Docker Container

hypervisor (VMM): hosts a computer, in software, on which SW runs.

- Type-1 (Hyper-V, Xen)
- Type-2 (VirtualBox, VMware)

library OS: OS as user-mode library. “Multiple OS” on a shared system.

- Graphene

containerization: restricted execution environment.

- `chroot, docker, lxc`

for completeness; sys-admins don't really ever do this

isolation done where:
inside the program

Intra-Application

code-rewriting: rewrite code to introduce isolation into it.

monitors, binary instrumentation

compiler: program analysis rejects program that do not have isolation.

Java memory safety, CompCert IFC languages

system-loading: force system to use customized libs that do access control

`boxify`

Summary

Summary

Arms Race

now you know some of what **Security Analysts**, **Forensic Analysts**, and **System Administrators** do.

- detect vulnerabilities
  tools: *scan*, look up *CVE*
- audit for attacks
  how/who/when/where
  - manual: visualization
  - automatic: report/act
- security w/o building security in

adblock,
fuckadblock
fuckfuckadblock,
...

important: limits; e.g. pattern-based approaches can be circumvented